The transportation problem and the diameters of transportation polytopes

Edward D. Kim

University of Wisconsin-La Crosse

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Slides at http://websites.uwlax.edu/ekim/talks/

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Example: Linear Program

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subject to $3x + 2y + z \le 10$
 $2x + 5y + 3z \le 15$
 $x, y, z \ge 0.$

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The Hirsch Conjecture

 $\Delta(d,f)$ = maximum diameter among d-polytope with f facets



The Hirsch Conjecture

 $\Delta(d, f)$ = maximum diameter among *d*-polytope with *f* facets



Conjecture [Hirsch 1957] $\Delta(d,f) \leq f - d$

The Hirsch Conjecture

 $\Delta(d, f)$ = maximum diameter among *d*-polytope with *f* facets



Ex-Conjecture [Hirsch 1957] $\Delta(d,f) \leq f - d$

Theorem: Matschke, Weibel, Santos [2011]

$$\Delta(d,f) \geq \frac{21}{20}(f-d)$$
 for $d \geq 20$

suppliesdemandsAustin•DelhiBoston•Edmonton•Fukuoka

Charleston • Giza

supplies

Austin 8 •

Boston 8 •

demands

- 9 Delhi
- 1 Edmonton
- 3 Fukuoka
- Charleston 6 • 9 Giza



Linear Program: the transportation problem Given:

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Linear Program: the transportation problem Given:

the supply quantities



Linear Program: the transportation problem

Given:

the demand quantities



Linear Program: the transportation problem Given:

• it costs \$9/unit to ship from Austin to Delhi



Linear Program: the transportation problem Given:

• it costs \$2/unit to ship from Austin to Edmonton



Linear Program: the transportation problem Given:

• it costs \$7/unit to ship from Austin to Fukuoka



Linear Program: the transportation problem Given:

• it costs \$8/unit to ship from Austin to Giza



Linear Program: the transportation problem Given:

• it costs \$8/unit to ship from Boston to Delhi



Linear Program: the transportation problem Given:

• it costs \$3/unit to ship from Boston to Edmonton



Linear Program: the transportation problem Given:

• it costs \$9/unit to ship from Boston to Fukuoka



Linear Program: the transportation problem Given:

• it costs \$3/unit to ship from Boston to Giza



Linear Program: the transportation problem Given:

• it costs \$7/unit to ship from Charleston to Delhi



Linear Program: the transportation problem Given:

• it costs \$4/unit to ship from Charleston to Edmonton



Linear Program: the transportation problem

Given:

• it costs \$9/unit to ship from Charleston to Fukuoka



Linear Program: the transportation problem Given:

• it costs \$6/unit to ship from Charleston to Giza



Linear Program: the transportation problem

Given:

• supply/demand amounts and per-unit transportation costs, How much should be put on the plane from each supply to each demand to minimize the total cost of transporting?

Let $u \in \mathbb{R}^m$ and $v \in \mathbb{R}^n$ with positive entries.

Definition

$$\sum_{j=1}^{n} x_{i,j} = u_i \quad \forall i \qquad \text{and} \qquad \sum_{i=1}^{m} x_{i,j} = v_j \quad \forall j.$$

<i>x</i> _{1,1}	<i>x</i> _{1,2}	<i>x</i> _{1,3}	<i>x</i> _{1,4}	8
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The $m \times n$ transportation polytope *P* defined by

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 and $\sum_{i=1}^{m} x_{i,j} = v_j \ orall j$ and $x_{i,j} \ge 0 \ orall i, j$

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If the Hirsch Conjecture is true for transportation polytopes, then $diam(P) \le f - d \le m + n - 1$.

Example: m = 3, n = 4

If $u = (8, 8, 6) \in \mathbb{R}^m$ and $v = (9, 1, 3, 9) \in \mathbb{R}^n$,

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If $u = (8, 8, 6) \in \mathbb{R}^m$ and $v = (9, 1, 3, 9) \in \mathbb{R}^n$, then the polytope *P* contains (among others), the following points:



Both the point on the left and on the right happen to be vertices due to a Theorem of Klee and Witzgall.

Same slice contains the same single support entry





Vertices X' and Y' have

Same slice contains the same single support entry





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Vertices X' and Y' have

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Quadratic Bound

Theorem: van den Heuvel, Stougie [2002]

Every $m \times n$ transportation polytope has diameter at most $(m + n)^2$.

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Lemma

Given two arbitrary vertices *X* and *Y* of an $m \times n$ transportation polytope *P*, there are vertices *X'* and *Y'* of *P* such that:

• The same slice in X' and Y' contains the same single support entry

•
$$\operatorname{dist}_{P}(X, X') + \operatorname{dist}_{P}(Y, Y') \le 2(m + n - 2).$$















Linear Bound

Theorem: Brightwell, van den Heuvel, Stougie [2006] Every $m \times n$ transportation polytope has diameter at most 8(m + n - 1).

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$$\operatorname{dist}_P(X, X') + \operatorname{dist}_P(Y, Y') \leq 8.$$





Better Linear Bound

Theorem: Hurkens [2009]

Every $m \times n$ transportation polytope has diameter at most 4(m + n - 1).

Better Linear Bound

Theorem: Hurkens [2009]

Every $m \times n$ transportation polytope has diameter at most 4(m + n - 1).

Lemma

Given two arbitrary vertices *X* and *Y* of an $m \times n$ transportation polytope *P*, there is an integer k > 0 such that

- X' and Y' are vertices of P
- The same *k* slices in *X'* and *Y'* each contain the same single support entry
- dist_P(X, X') + dist_P(Y, Y') $\leq 4k$.

$2 \times n$ Transportation Polytopes

Theorem: De Loera, K. [2014]

Every $2 \times n$ transportation polytope satisfies the Hirsch Conjecture.

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Theorem: De Loera, K. [2014]

Every $2 \times n$ transportation polytope satisfies the Hirsch Conjecture.

Lemma

Every vertex *X* of a non-degenerate $2 \times n$ transportation polytope contains a unique column with two strictly positive coordinates.



3-way Transportation Polytopes

Definition Given $u \in \mathbb{R}^m$, $v \in \mathbb{R}^n$ and $w \in \mathbb{R}^p$,



3-way Transportation Polytopes

Definition

Given $u \in \mathbb{R}^m$, $v \in \mathbb{R}^n$ and $w \in \mathbb{R}^p$, the 3-way transportation polytope given by 1-marginals is defined in *mnp* non-negative variables $x_{i,j,k} \in \mathbb{R}_{\geq 0}$ with the m + n + p equations



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$$\sum_{j,k} x_{i,j,k} = u_i \quad \forall i,$$


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Significance of 3-way Transportation Polytopes



Theorem: De Loera, Onn [2006]

Given any rational polytope P,

Significance of 3-way Transportation Polytopes



Theorem: De Loera, Onn [2006]

Given any rational polytope P,

• There is a 3-way transportation polytope *Q* given by 1-marginals with a face that is isomorphic to *P*.

Significance of 3-way Transportation Polytopes



Theorem: De Loera, Onn [2006]

Given any rational polytope P,

- There is a 3-way transportation polytope *Q* given by 1-marginals with a face that is isomorphic to *P*.
- Moreover, the polytope *Q* can be computed in polynomial time (in the description of *P*).

















Quadratic Bound for Axial Transportation Polytopes

Theorem: De Loera, K., Onn, Santos [2009]

Every 3-way axial $m \times n \times p$ transportation polytope has diameter at most $2(m + n + p)^2$.

Quadratic Bound for Axial Transportation Polytopes

Theorem: De Loera, K., Onn, Santos [2009]

Every 3-way axial $m \times n \times p$ transportation polytope has diameter at most $2(m + n + p)^2$.

Lemma

Given two arbitrary vertices *X* and *Y* of an $m \times n \times p$ axial transportation polytope *P*, there are vertices *X'* and *Y'* of *P* such that:

- The same slice in X' and Y' contains the same single support entry
- $\operatorname{dist}_{P}(X, X') + \operatorname{dist}_{P}(Y, Y') \le 4(m + n + p 1).$

















Thank you!

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